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DESI	GNATED/ELECTED OFFICE (DO/EO/US) CERNING A FILING UNDER 35.U.S.C. 371	ATTORNEY'S DOCKET NO. A34394 PCT U	SA
		U.S. APPLICATION NO. 86	87 97
INTERNATIONAL APPLICATION NO. PCT/DE99/04050	INTERNATIONAL FILING DATE 21 DECEMBER 1999	PRIORITY DATE CLAIMED 22 DECEMBER 19	98
THE OF INVENTION DEVICE AND METHO	D FOR GENERATING AND EXECUTING CO Y LONG INSTRUCTION WORD PROCESS	OMPRESSED OR	
APPLICANT(S) FOR DO/EO/US Matthias Weiss			
1. This is a FIRST submission of item 2. [] This is a SECOND or SUBSEQUE 3. [] This express request to begin natio until the expiration of the applicable time 4. [] A proper Demand for International Application a. [] is transmitted herewith (required as [] is transmitted by the c. [] is not required, as the application of the International Action of the amendments to the second of the Internation of the invention of the invention of the invention of the invention of the International Preliminary action of the International Preliminary action of the Internation of Supplement Excopy of International Application of Supplement Excopy of Internat	International Bureau. International Bureau. International Bureau. International Bureau. International Bureau. International English (35 U.S.C. 371(c)(2)). Report (PCT/ISA/210) International Bureau International Bureau International Bureau International Bureau International Bureau International PCT Article 19 (35 U.S.C. 37 (c)(4)). International Preliminary Examination Report under Examination Report under Examination Report (PCT/IPEA/409) International Preliminary Examination Report under Examination Report (PCT/IPEA/409) International Preliminary Examination Report (PCT/IPEA/409)	er 35 U.S.C. 371. It any time rather than de 2 and 39(I). Imonth from the earliest of Bureau). Iffice (RO/US). I Bureau). as NOT expired. I(c)(3)). Inder PCT Article 36 (35) I 37 CFR 3.28 and 3.31	elay examination claimed priority date. 5 U.S.C. 371(c)(5)).

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17. The following fees are submitted:	<u>-</u>			_C/	ALCULATIONS	PTO USE ONLY
Basic National Fee (37 CFR 1.492(a))						
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Nor international search fee (37 CFR 1.44: Report not prepared by the EPO or JPO (1	5(a)(2)) paid to	USPTO and Internation	onal Search			
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SEND ALL CORRESPONDENCE TO: Ronald B. Hildreth		1 mos	at H	_	educh	
BAKER BOTTS L.L.P.		Attorney: Ronald	B. Hild eth		PT	O Reg: 19,498
30 Rockefeller Plaza			▼	2	21 JUNE 2001	
New York, New York 10112-4498					Date	

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BAKER BOTTS LLP

Attorney Docket Number: A34394 PCT USA

Title:

DEVICE AND METHOD FOR GENERATING AND EXECUTING COMPRESSED PROGRAM OF A

VERY LONG INSTRUCTION WORD PROCESSOR

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Rec'd PCT/PTO 25 SEP 2001

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PTO is the application which the inventor executed by signing the oath or declaration.

Also	enclose	ed are:	
1. (eck in the amount of \$130 - Other than Small Entity \$ 65 - Small Entity	\$_130
		representing payment of the surcharge due for late filing of the Declaration pursuant to 37 C.F.R. 1.492(e);	
2.	0	An extension of time to file the Response to Missing Parts is respectfully requested. The required fee, calculated pursuant to 1.136(a), is enclosed in the amount of	\$
3.	()	Other	\$
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Respectfully submitted,

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NY02:348218.1

09/868797 JC18 Rec'd PCT/PTO 2 1 JUN 2001 A34394 PCT USA

PATENT

IN THE UNITED STATES PATENT AND TRADEMARK OFFICE

Applicant

Matthias Weiss

Serial No.

To be Assigned

Filed

To be Assigned

For

DEVICE AND METHOD FOR GENERATING AND EXECUTING

COMPRESSED PROGRAM OF A VERY LONG INSTRUCTION

WORLD PROCESSOR

PRELIMINARY AMENDMENT

Assistant Commissioner of Patents

Washington, D.C. 20231

Sir:

Preliminary to the examination of the above-identified application, please make the following amendment to the claims:

In the Claims:

Cancel claims 1 to 10.

Add the following new claims 11 to 20.

11. A method for controlling functional units in a processor in a configuration, in which a sequence of primary instruction words consisting of multiple instruction word parts and originating from a translation of a program code is compressed and stored as a sequence of related program words, and according to which, in a subsequent execution phase, sequential secondary instruction words consisting of a plurality of instruction word parts and having the full instruction word width needed to control all functional units are generated from the sequence of program words, comprising:

as a result of the configuration, a program word has a first characteristic of a primary instruction word from a first group of preceding primary instruction words, which has the greatest similarity to the primary instruction word associated with the program word, and contains instruction word parts that differentiate the primary instruction word belonging to the program word from the primary instruction word belonging to the first characteristic;

in the execution phase, storing each instruction word in a second group of secondary instruction words corresponding in number to the first group is provided with a second characteristic;

in accordance with the first characteristic contained in the program word, ascertaining a secondary instruction word corresponding to the associated primary instruction word from the second group via the associated second characteristic; and

generating the secondary instruction word corresponding to the program word such that the instruction word parts contained in the program word are exchanged in the secondary instruction word from the second group.

12. The method in accordance with claim 11, wherein:

the first group consists of a first number of primary instruction words that directly precede the primary instruction word in question; and

the second group consists of a second number of secondary instruction words that is at least equal to the first number, where, prior to the generation of the next sequential secondary instruction word, each most recent secondary

instruction word is appended to the second group as the last word, and the first secondary instruction word to have been added and that is in excess of the second number is removed from the second group.

- 13. The method in accordance with claim 11, wherein the newly generated secondary instruction word is appended to the second group in that the former is stored in place of the secondary instruction word that was used for its generation.
- 14. The method in accordance with claim 11, wherein the newly generated secondary instruction word is not stored.
- 15. The method in accordance with 11, wherein the first characteristic is formed as a minimum code distance between the primary instruction word belonging to the program word in question and the primary instruction word with the greatest similarity.
- 16. The method in accordance with 11, wherein the second characteristic consists of an address corresponding to the first characteristic that is the address of a preceding secondary instruction word in a memory used for storage of the second group.
- 17. The method in accordance with claim 11, wherein the program word consists of a number of instruction word parts that corresponds to the number of instruction

word parts to be differentiated that occurs most frequently within the configuration, and in that a plurality of program words are used to assemble secondary instruction words that require more than the number of instruction words stored in one program word for the secondary instruction word used for generation.

- 18. The method in accordance with claim 11, wherein the instruction word parts are compressed in one program word by reducing the bit width to the extent that it is possible to represent the most frequently occurring instruction word parts; and in that multiple program words are used when instruction word parts occur that require a greater bit width in order to be represented.
- 19. The method in accordance with claim 18, wherein the width of the instruction word parts in the program word is halved, and one or two program words are provided for representation of the instruction word parts.
- 20. A processor arrangement for carrying out the method of claim 1, comprising:

 a plurality of functional units;

 an instruction word memory associated with the functional units; and

 an instruction word buffer for storing instruction words that have already

 been generated and have a width that is at least the size of the bit width of the

 secondary instruction word; the instruction word buffer including a memory with

 selective line-by-line access.

R E M A R K S

This amendment eliminates multiple dependency in the claims and puts the claims in better U.S. format. No new matter is introduced by this amendment.

Respectfully submitted,

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[Translation from German]

WO 00/38049

PCT/DE99/04050

Device and Method for Generating and Executing Compressed Programs of a Very-Long-Instruction-Word Processor

Description

The invention relates to a method for controlling functional units in a processor, according to which, in a configuration, a sequence of primary instruction words consisting of multiple instruction word parts and originating from a translation of a program code is compressed and stored as a sequence of related program words. In a subsequent execution phase, sequential secondary instruction words consisting of a plurality of instruction word parts and having the full instruction word width needed to control all functional units are generated from the sequence of program words.

The invention also relates to a processor arrangement for carrying out the method having functional units, an instruction word memory associated with these functional units and an instruction word buffer for storing instruction words that have already been generated and have a width that is at least the size of the bit width of the secondary instruction word.

Processor arrangements of the aforementioned type contain functional units that operate in parallel to one another and that are controlled at every clock cycle by an instruction word. The particular instruction word is extracted from a program word that is taken from a program memory.

For their part, the instruction words consist of a plurality of instruction word parts, where each individual instruction word part serves to control one functional unit.

To improve the performance of processor arrangements, the goal is an increase in the processing width, which makes it necessary to increase the number of functional units. In general, this increases the bit width of the instruction words and thus also of the program words. The consequence of this is the provision of corresponding storage space in the program memory, which occupies the majority of the area on the semiconductor chip.

Since the size of the program words determines the size of the program memory, the goal is to reduce the size of the program words in order to reduce the need for memory space. A number of compression methods for this purpose are known. The most obvious method is described in the report by H. Weiss and G. Fettweis, [in English:] "Dynamic Codewidth Reduction for VLIW Instruction Set Architectures in Digital Signal Processors" (Proceedings of the 3rd International Workshop on Signal and Image Processing IWSIP '96, pages 517 to 520). In this method, the program words are assembled from sequential primary instruction words in such a way that secondary instruction words can be subsequently reproduced therefrom in that a secondary instruction word (VLIW),

once it has been created, is written to an instruction word memory, and, in order to produce the next secondary instruction word, only those instruction word parts in the stored secondary instruction word are exchanged which differ between the stored secondary instruction word and the secondary instruction word to be generated. Consequently, the program word need only contain the information specifying which instruction word part differs and with what content it differs. It is thus possible to design the program words to be very narrow and thus save memory space.

However, when there are great differences between the stored secondary instruction word and the secondary instruction word to be created, the width of the program word must be increased if these relatively great differences occur frequently, which entails the disadvantage of a relatively large memory space, or else the differences must be distributed over multiple program words. Thus, the secondary instruction word must be created from multiple program words over multiple clock cycles. This results in the disadvantage that it requires a relatively long time.

Hence, the object of the invention is to increase operating speed in an application-specific manner while retaining a small program word width.

As regards the method, the object is attained in that, as the result of the configuration a program word contains a first characteristic of a primary instruction word from a first group of preceding primary instruction words which has the greatest degree of correspondence with the primary instruction word associated with the program word, and contains instruction word parts which

differentiate the primary instruction word belonging to the program word from the primary instruction word belonging to the characteristic. In the execution phase, a second group — corresponding in number to the first group — of secondary instruction words, each of which is equipped with a second characteristic, is stored. By means of the first characteristic contained in the program word, a secondary instruction word corresponding to the associated primary instruction word is ascertained from the second group via the associated second characteristic, and the secondary instruction word corresponding to the program word is generated in that the instruction word parts contained in the program word are exchanged in the secondary instruction word from the secondary group.

As early as the configuration phase, the program words can be constructed such that they contain only the minimum possible information for later generation of a secondary instruction word in the execution phase. This is achieved by referring back to the preceding primary instruction words that have the greatest correspondence with the primary instruction word to be configured. Since the secondary instruction words are to be generated during the execution phase with the same width and in the same sequence as the primary instruction words, the execution phase proceeds essentially in the reverse order to the configuration phase, and the already generated secondary instruction words corresponding to the primary instruction word most similar in the configuration phase to a current primary instruction word are used one at a time to generate a new secondary instruction word. Since the instruction word with the greatest similarity or the greatest correspondence is always referred back to, the amount

of information necessary to generate a new secondary instruction word can be kept as small as possible. In this way, it is possible to minimize the storage requirement of a program memory.

In one embodiment of the method, provision is made for the first group to consist of a first number of primary instruction words directly preceding the primary instruction word in question. The second group consists of a second number of secondary instruction words that is at least equal to the first number, where, prior to the generation of the next sequential secondary instruction word, the most recent secondary instruction word is appended to the group as the last word, and the first secondary instruction word to have been added that is in excess of the second number is removed from the second group. The group thus always consists of the immediately preceding instruction words, one of which has a greatest possible similarity when the instruction words within a group differ in content.

Another possible embodiment of the method is for the newly generated secondary instruction word to be appended to the second group in that this word is stored in place of the secondary instruction word that was used to generate it.

This variant of the process precludes the possibility of the group filling up with instruction words having the same content when the same function is passed through multiple times.

Another possibility is that the newly generated secondary instruction word is not stored. In this case, the same set of previously generated instruction words is always accessed each time a secondary instruction word is generated,

which is advantageous when these stored instruction words are suitable as pattern words for the generation of other secondary instruction words, with the result that storage processes can be avoided by this means.

In one embodiment of the method, provision is made for the first characteristic to be formed as a minimum code distance between the primary instruction word belonging to the program word in question and the primary instruction word with the greatest similarity.

Another embodiment of the method provides for the second characteristic to consist of an address corresponding to the first characteristic that is the address of a preceding secondary instruction word in a memory used for storage of the second group.

Both of these embodiments exhibit an especially simple generation of features.

A reduction in the program word width, and thus an associated reduction in the width of the program memory, can be achieved in that the program word consists of a number of instruction word parts that corresponds to the number of instruction word parts to be differentiated which occurs most frequently within the configuration, and in that a plurality of program words are used to assemble secondary instruction words that require more than the number of instruction words stored in one program word for the secondary instruction word used for generation. Consequently, the width of the program memory satisfies the most common application cases, and no delay arises in the generation of secondary instruction words. Only in the relatively few cases where the number of

instruction word parts to be changed exceeds the width of a program word are two or more program words required in two or more clock cycles to produce the secondary instruction word.

A further reduction in program word width can be achieved in that the instruction word parts are compressed in one program word. This is accomplished in that the bit width is reduced to the extent that it is possible to represent the most frequently occurring instruction word parts, and in that multiple program words are used when instruction word parts occur that require a greater bit width in order to be represented.

For this type of compression of the program word, it is especially useful for the width of the instruction word parts in the program word to be halved and for one or two program words to be provided for representation of the instruction word parts. For example, if the program word normally has a width of 8 bits, 256 combinations are available. However, a large number of these combinations is required only extremely rarely or not at all. Accordingly, the program word can be provided with a width of 4 bits, which is sufficient to represent the 16 most frequently occurring combinations. If a combination other than these should need to be represented, two or more program words are used for this purpose. However, since this occurs only very rarely, the savings in storage space for the program memory that is achieved is greater than the possible expenditure of time for generation of rare combinations.

The object is also achieved in accordance with the invention by a processor arrangement that is characterized in that the instruction word buffer

consists of a memory with selective line-by-line access. In contrast to the prior art shown, this affords the possibility of directly accessing various stored instruction words in order to produce a new secondary instruction word. Thus, it is not necessary to use just any stored instruction word for generating a new secondary instruction word, but rather the particular stored instruction word can be used that has the greatest similarity to the secondary instruction word to be generated. Hence, the resource requirements for changes are relatively small, requiring only a small bit width for the program word, and thus a small program memory.

The invention is described in detail below with reference to an example embodiment. The appertaining drawings show:

- Fig. 1 a flowchart of the inventive method, and
- Fig. 2 the principle of the inventive generation of the primary and secondary instruction words.

In the method represented in Fig. 1 for control of functional units 1 in a processor 2, a sequence of primary instruction words 5 is generated from a plurality of instruction word parts 6 from a program code 3 by means of a translation 4 in a configuration phase. This sequence of primary instruction words 5 is compressed and stored in a program memory 8 as a sequence of related program words 7.

The program words 7 consist of a number of instruction word parts 6 corresponding to the most frequently occurring number of instruction words 6 to

be differentiated within the configuration. In order to compose secondary instruction words 9, more than the number of instruction word parts 6 stored in one program word 7 may be needed. Then a plurality of program words 7 are used to do so.

In a subsequent execution phase, sequential secondary instruction words 9 consisting of a plurality of instruction word parts 6 and having the full instruction word width needed to control all functional units 1 are generated from the sequence of program words 7.

In the result of the configuration, a program word 7 has a first characteristic 10 of a primary instruction word 5 from a first group 11 of preceding primary instruction words 5, which has the greatest similarity to the primary instruction word 5 associated with the program word 7, and contains instruction word parts 6 that differentiate the primary instruction word 5 belonging to the program word 7 from the primary instruction word 5 belonging to the characteristic. The first characteristic 10 is formed as a minimum code distance between the primary instruction word 5 belonging to the relevant program word 7 and the primary instruction word 5 with the greatest similarity.

In the execution phase, a second group 12 of secondary instruction words 9 — corresponding in number to the first group 11 — each of which is provided with a second characteristic 13, is stored. The second characteristic 13 is formed from a line number of a memory 14 that serves to store the second group 12.

By means of the first characteristic 10 contained in the program word 7, a secondary instruction word 9 corresponding to the associated primary instruction word 5 is ascertained from the second group 12 via the associated second characteristic 13. The secondary instruction word 9 corresponding to the program word is generated in that the instruction word parts 6 contained in the program word 7 are exchanged in the secondary instruction word 9 from the second group 12.

The first group 11 consists of a first number of the primary instruction words 5 preceding the primary instruction word 5 in question. The second group 12 consists of a second number of secondary instruction words 9, which is at least equal to the first number, where, prior to the generation of the next sequential secondary instruction word 9, each most recent secondary instruction word 9 can be appended to the second group 12 as the last word. The first secondary instruction word 9 to have been added that is in excess of the second number is removed from the second group 12.

Method for controlling functional units in a processor and processor arrangement for carrying out the method

List of Symbols

1	functional unit
2	processor
3	program code
4	translation
5	primary instruction word
6	instruction word part
7	program word
8	program memory
9	secondary instruction word
10	first characteristic
11	first group
12	second group
13	second characteristic
14	memory

<u>Claims</u>

Method for controlling functional units in a processor, according to which, in a configuration, a sequence of primary instruction words consisting of multiple instruction word parts and originating from a translation of a program code is compressed and stored as a sequence of related program words, and according to which, in a subsequent execution phase, sequential secondary instruction words consisting of a plurality of instruction word parts and having the full instruction word width needed to control all functional units are generated from the sequence of program words characterized in that,

in the result of the configuration, a program word (7) has a first characteristic (10) of a primary instruction word (5) from a first group (11) of preceding primary instruction words (5), which has the greatest similarity to the primary instruction word (5) associated with the program word (7), and contains instruction word parts (6) that differentiate the primary instruction word (5) belonging to the program word (7) from the primary instruction word (5) belonging to the first characteristic (10)

and in that, in the execution phase, a second group (12) of secondary instruction words (9) — corresponding in number to the first group (11) — each of which is provided with a second characteristic (13), is stored, and

in that, by means of the first characteristic (10) contained in the program word (7), a secondary instruction word (9) corresponding to the associated primary instruction word (5) is ascertained from the second group (12) via the associated second characteristic (13), and the secondary instruction word (9) corresponding to the program word (7) is generated in that the instruction word parts (6) contained in the program word (7) are exchanged in the secondary instruction word (9) from the second group (12).

2. Method in accordance with claim 1, **characterized in that** the first group (11) consists of a first number of primary instruction words (5) that directly precede the primary instruction word (5) in question and

in that the second group (12) consists of a second number of secondary instruction words (9) that is at least equal to the first number, where, prior to the generation of the next sequential secondary instruction word (9), each most recent secondary instruction word (9) is appended to the second group (12) as the last word, and the first secondary instruction word (9) to have been added and that is in excess of the second number is removed from the second group (12).

3. Method in accordance with claim 1, **characterized in that** the newly generated secondary instruction word (9) is appended to the second

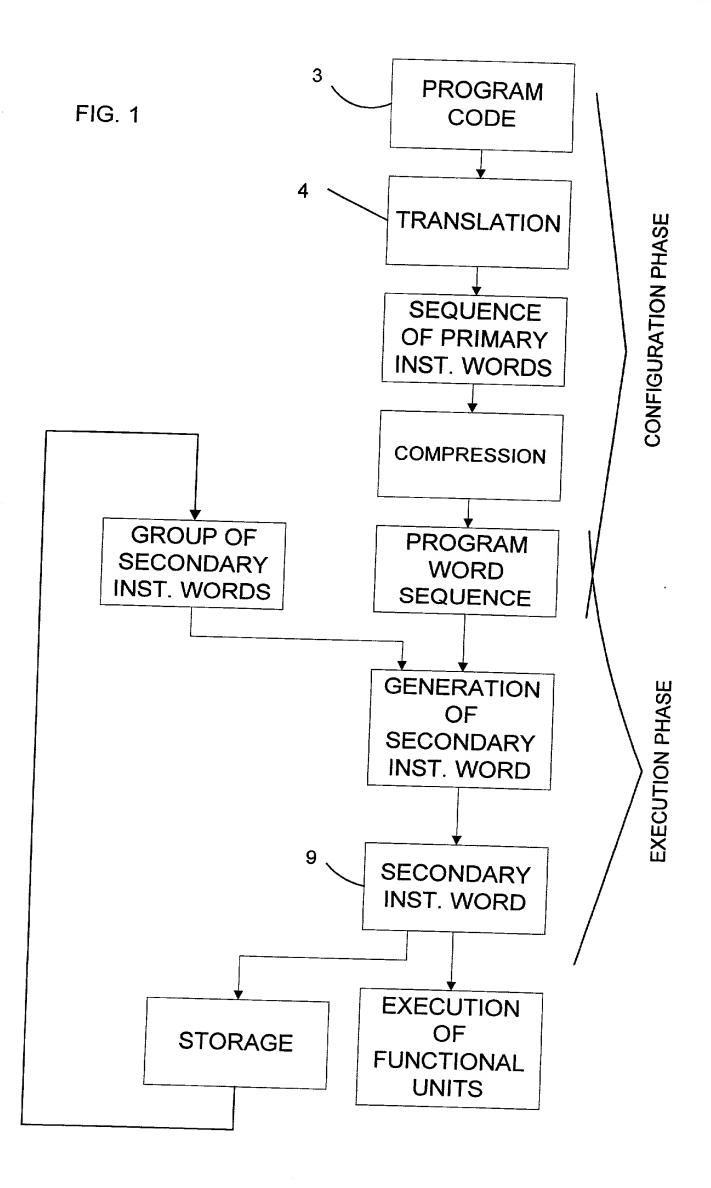
group (12) in that the former is stored in place of the secondary instruction word (9) that was used for its generation.

- 4. Method in accordance with claim 1, **characterized in that** the newly generated secondary instruction word (9) is not stored.
- 5. Method in accordance with one of claims 1 through 4, **characterized in that** the first characteristic (10) is formed as a minimum code distance
 between the primary instruction word belonging to the program word in
 question and the primary instruction word with the greatest similarity.
- 6. Method in accordance with one of claims 1 through 5, **characterized in that** the second characteristic (13) consists of an address corresponding
 to the first characteristic that is the address of a preceding secondary
 instruction word in a memory (14) used for storage of the second group
 (12).
- 7. Method in accordance with one of claims 1 through 6, **characterized in that** the program word (7) consists of a number of instruction word parts
 (6) that corresponds to the number of instruction word parts (6) to be
 differentiated that occurs most frequently within the configuration, and in
 that a plurality of program words (7) are used to assemble secondary
 instruction words (9) that require more than the number of instruction

- words (9) stored in one program word (7) for the secondary instruction word (9) used for generation.
- 8. Method in accordance with one of claims 1 through 7, **characterized in that** the instruction word parts (6) are compressed in one program word
 (7) by reducing the bit width to the extent that it is possible to represent
 the most frequently occurring instruction word parts (6), and in that
 multiple program words (7) are used when instruction word parts (6) occur
 that require a greater bit width in order to be represented.
- 9. Method in accordance with claim 8, **characterized in that** the width of the instruction word parts (6) in the program word (7) is halved, and one or two program words (7) are provided for representation of the instruction word parts (6).
- 10. Processor arrangement for carrying out the method, having functional units, an instruction word memory associated with these functional units and an instruction word buffer for storing instruction words that have already been generated and have a width that is at least the size of the bit width of the secondary instruction word, **characterized in that** the instruction word buffer consists of a memory (14) with selective line-by-line access.

ABSTRACT OF THE DISCLOSURE

The invention relates to a method for controlling functional units in a processor, according to which in a configuration a sequence of primary instruction words which consists of several instruction word parts and originates from a translation of a program code is compressed and stored as a sequence of associated program words. The invention also relates to a processor system for carrying out this method. The aim of the invention is to increase operating speed in an application-specific manner while retaining a low program word width. To this end, as regards the method, a program word contains a first characteristic of a primary instruction word and instruction word parts which differentiate the primary instruction word belonging to the program word from the primary instruction word belonging to the characteristic. By means of the first characteristic contained in the program word a secondary instruction word is generated by exchanging the instruction word parts contained in the program word with those in a preceding secondary instruction word. On the system side the aim of the invention is solved by providing for the instruction word buffer to consists of a memory with optional line-by-line access.



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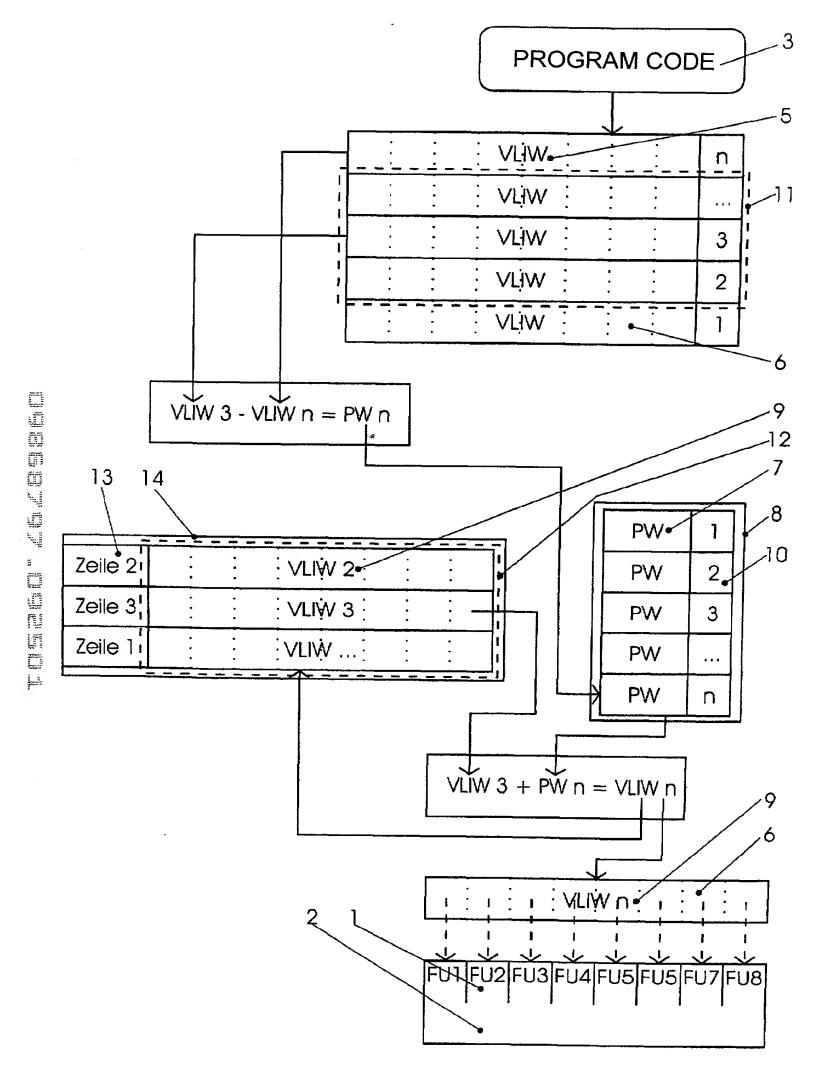


Fig. 2

BAREKBOTTS L L.P. FILE NO.: A34394 PCT USA

COMBINED DECLARATION AND POWER OF ATTORNEY

(Original, Design, National Stage of PCT, Divisional, Continuation or C-I-P Application)

As a below named inventor, I hereby declare that: Matthias Weiss

(d) [] no such applications have been filed.

(e) [x] such applications have been filed as follows:

My residence, post office address and citizenship are as stated below next to my name; I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled:

DEVICE AND METHOD FOR GENERATING AND EXECUTING COMPRESSED PROGRAMS OF A VERY LONG INSTRUCTION WORD PROCESSOR

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(a)	[] is atta	ached hereto.
(b)	[X] was	s filed on Jun 21, 2001 as Application Serial No09/868,797 and was amended on (if
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		described and claimed in PCT International Application No. filed on and was amended on (if
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matter		vledge the duty to disclose information which is material to the patentability of the subject in this application in accordance with Title 37, Code of Federal Regulations § 1.56.
[]]	n compli	ance with this duty there is attached an information disclosure statement. 37 CFR 1.98.
	w. 1	Priority Claim
4.	I hereby	claim foreign priority benefits under Title 35, United States Code, § 119(a)-(d) of any foreign
applica	ition(s) it	or patent or inventor's certificate or of any PCT International Application(s) designating at least
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FILE NO.: A34394 PCT USA

PRIOR FOREIGN/PC APPLICATION	T APPLICATION(S) FILED WIT	THIN 12 MONTHS (6 MC	ONTIIS FOR DESIGN) PE	ZIOR TO SAID
COUNTRY	APPLICATION NO.	DATE OF FILING (day, month, year)		PRIORITY CLAIMED UNDER 35 USC 119
				[] YES NO[]
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ALL FOREIGN APPL APPLICATION	JCATION[S], IF ANY, FILED M	ORE THAN 12 MONTH	S (6 MONTHS FOR DESI	GN) PRIOR TO SAID
Germany	198 59 389.9	22 December 1998		[x] YES NO []
International	PCT/DE99/04050	21 December 1999		[x] YES NO []
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(Application S	Serial No.)	(Filing Date)	(patented, pen	ding, abandoned)
17,021; Francis J. 18,839; Ronald B. No. 24,316; Richar No. 27,551; James 29,705; Robert C. S 32,439; Rochelle K 32,689; Lisa B. Kolof the firm of BAK	Proof, I hereby appoint Dana I Hone, Rcg. No. 18,662; Jo Hildreth, Reg. No. 19,498; d G. Berkley, Reg. No. 26,946; J. Maune, Reg. No. 26,946; Scheinfeld, Reg. No. 31,300; K. Seide Reg. No. 32,300; G le, Reg. No. 35,225; Neil P. ER BOTTS L.L.P., with of the this application and to transport the series of the series of the proof of the series of the se	seph D. Garon, Reg. Thomas R. Nesbitt, J. 65; Richard S. Clark 5; John D. Murnane, John A. Fogarty, Jr., Jary M. Butter, Reg. No. 38,3 fices at 30 Rockefells	No. 20,420; Arthur r., Reg. No. 22,075; l., Reg. No. 26,154; Br Rcg. No. 29,836; He Rcg. No. 22,348; Loui No. 33,841; Marta E. I 106; and Paul A. Raguer Plaza. New York.	S. Tenser, Reg. No Robert Neuner, Reg adley B. Geist, Reg enry Tang, Reg. No s S. Sorell, Reg. No Delsignore, Reg. No isa, Reg. No. 38,58 New York 10112

SEND CORRESPONDENCE TO: BAKER BOTTS L.L.P. 30 ROCKEFELLER PLAZA, NEW YORK, N.Y. 10112 CUSTOMER NUMBER: 21003	DIRECT TELEPHONE CALLS TO: BAKER BOTTS L.L.P. (212) 408-2500

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

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